

嵌入式系統

單元6：uC/OS-II為例之RTOS

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單元學習目標與大綱

- uC/OS-II簡介
- uC/OS-II排程
- uC/OS-II軟體架構
- uC/OS-II Task同步-Semaphore
- uC/OS-II 訊息傳遞-Mailbox
- uC/OS-II 多訊息傳遞-Message Queue
- uC/OS-II 多事件觸發-Flags
- uC/OS-II 記憶體管理
- uC/OS-II 時間管理

Part 1

uC/OS-II簡介

uC/OS-II (1)

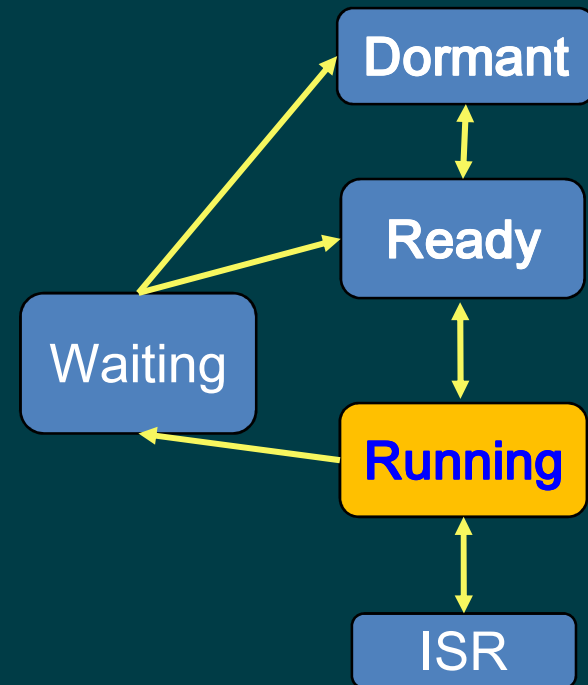
- **Real-time kernel**
- **Priority-based preemptive** 排程機制
 - 目前在即行的Task即為ready-list內**最高優先權的Task**
- 基本上有**64個priority number**
 - **Number愈小，優先權愈大**
- **每個task都要指定一個priority number**
 - 不同task，**不可設定相同的priority number**
 - **Priority number 0, 1, 2, 3為系統保留的號碼，不要使用**

uC/OS-II (2)

- 兩個system tasks: **Idle Task, Statistics Task**
 - 對應的priority number 為 **63與62**
- Priority number即為**task**之ID
- 相關資訊可參加書籍 **MicroC/OS-II: The Real-Time Kernel** (A complete portable, ROMable scalable preemptive RTOS), Jean J. LaBrosse, CMP Books

uC/OS-II Task States

- 5種可能的狀態
 - **Dormant state**: 被create但OS還不知其存在
 - **Ready state**: 在ready-list等待被執行
 - **Running state**: 目前正在CPU執行
 - **Waiting state**: 在waiting-list等待事件發生
 - **ISR state**: 在執行時，被中斷插斷



uC/OS-II Task

- 為一無窮迴圈的函數
- 建立Task
 - OSTaskCreate()
 - OSTaskCreateExt()
- 刪除Task
 - OSTaskDel()
- 暫停Task
 - OSTaskSuspend()
- 恢復Task
 - OSTaskResume()

無窮迴圈

Task name

```
void YourTask (void *pdata)
{
  for (;;) {
    /* USER CODE */
    Call one of uC/OS-II's services:
    {
      OSMboxPend();
      OSQPend();
      OSSemPend();
      OSTaskDel(OS_PRIO_SELF);
      OSTaskSuspend(OS_PRIO_SELF);
      OSTimeDly();
      OSTimeDlyHMSM();
    }
    /* USER CODE */
  }
}
```

uC/OS-II Critical Section保護

- 利用 **disable interrupts** 方式進入 **critical sections**，離開時再 **reenable interrupts**
- uC/OS-II 定義兩個 **macros** 可對中斷 **disable/enable**
 - **OS_ENTER_CRITICAL()**
 - **OS_EXIT_CRITICAL()**
- 三種方式 **實現 macros**，主要是採用方法二
 - **#define OS_ENTER_CRITICAL () **
 **asm(" PUSH PSW") **
 asm("DI")
 - **#define OS_EXIT_CRITICAL () **
 asm("POP PSW")

Part 2


uC/OS-II编程

uC/OS-II排程技巧

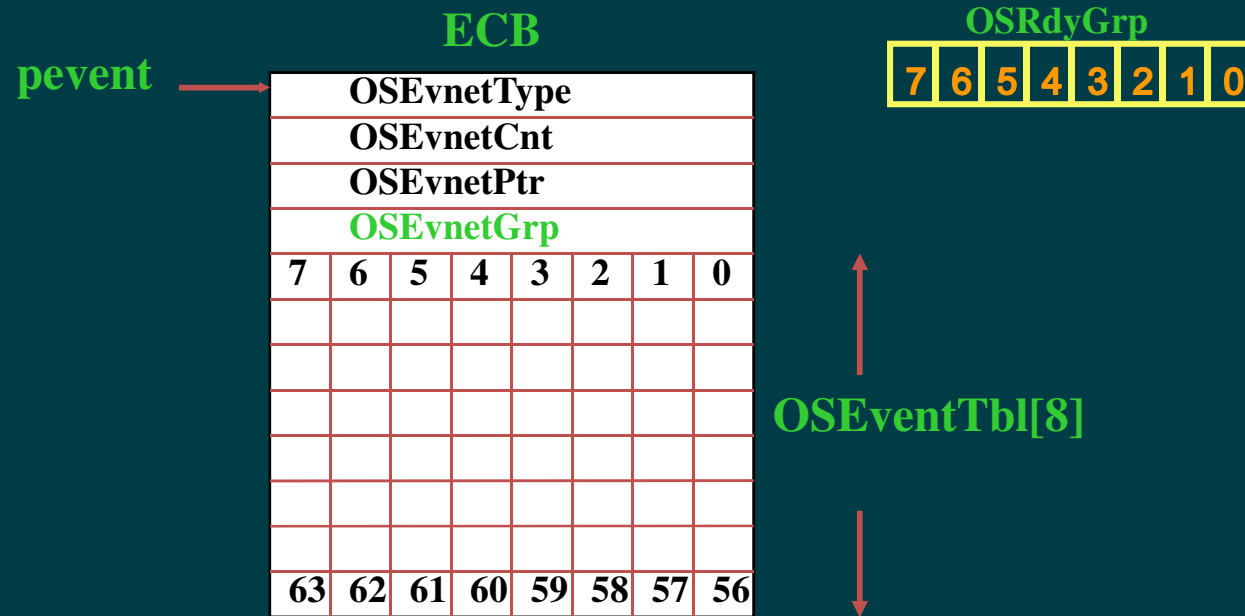
- 每個 semaphore, mailbox, and queue 都有自己的ECB (Event Control Block)
 - 當有Task等待semaphore, mailbox,或 queue, 即將此等待Task ID放在對應的ECB
 - 當等待事件滿足時, 即藉由ECB找出所有等待Task中, 最高優先權的Task, 將其由waiting state改為ready state

ECB

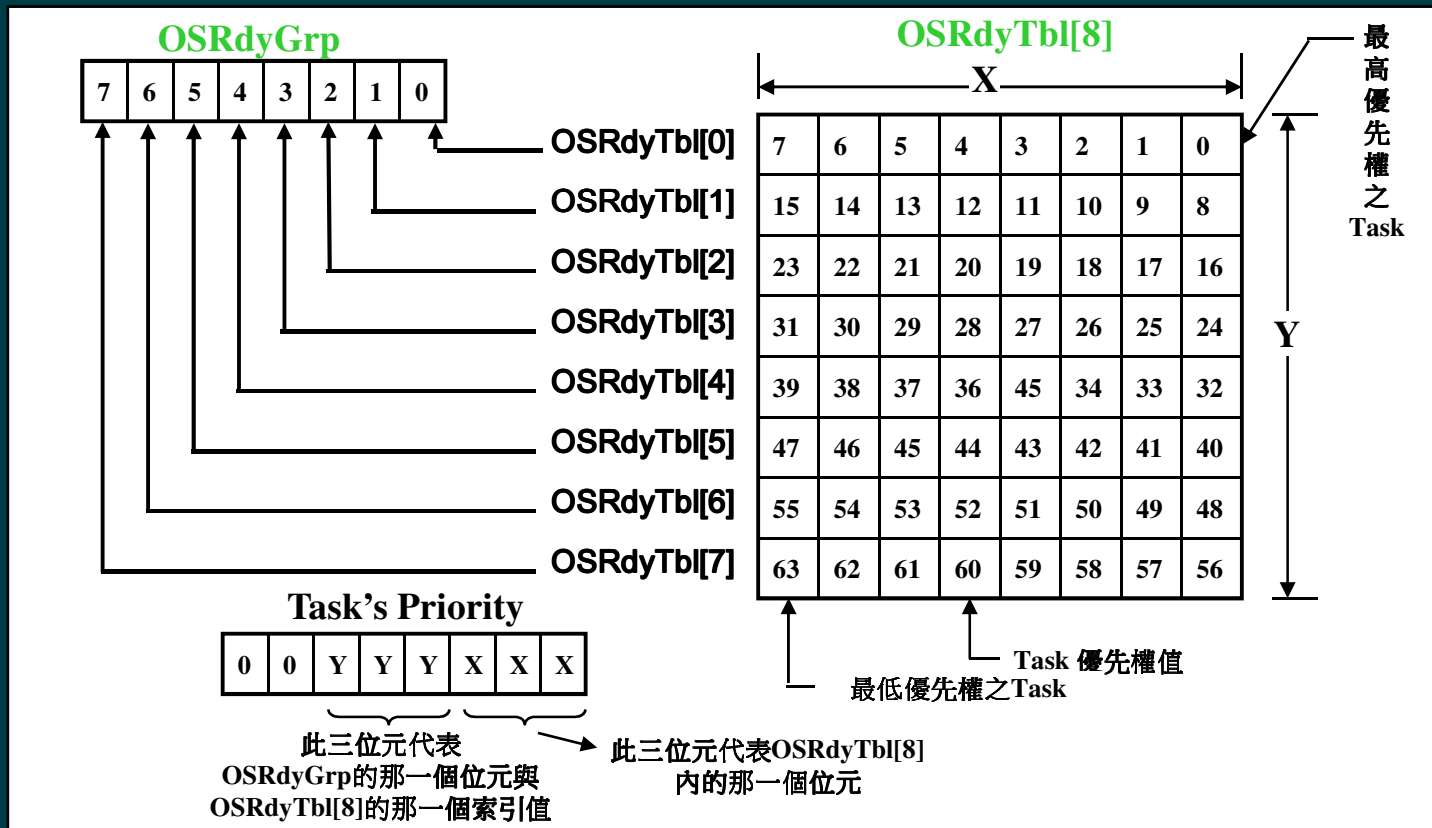
```
typedef struct {  
    INT8U  OSEventType;          /* Event type          */  
    INT8U  OSEventGrp;          /* Group for wait list */  
    INT16U OSEventCnt;          /* Count (when event is a semaphore) */  
    void  *OSEventPtr;          /* Ptr to message or queue structure */  
    INT8U  OSEventTbl[OS_EVENT_TBL_SIZE]; /* Wait list for event to occur */  
} OS_EVENT;
```



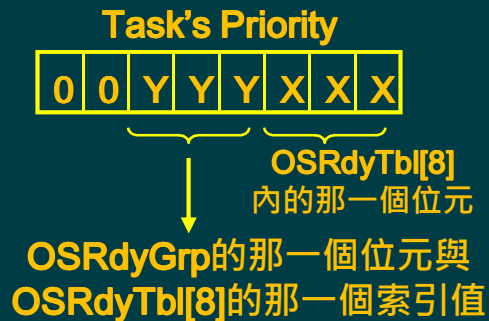
ECB主要是由EventGrp與EventTbl組成



EventGrp與EventTbl的關係



將Task加入Waiting List與移出



將Task加入Event的waiting list

```
pevent->OSEventGrp |= OSMapTbl[prio >> 3];  
pevent->OSEventTbl[prio >> 3] |= OSMapTbl[prio & 0x07];
```

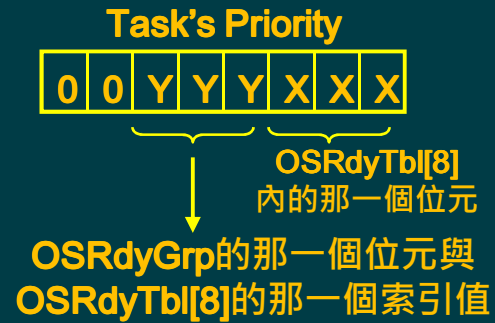
將Task由Event的waiting list移除

```
if ((pevent->OSEventTbl[prio >> 3] &= ~OSMapTbl[prio & 0x07]) == 0)  
{ pevent->OSEventGrp &= ~OSMapTbl[prio >> 3]; }
```

OSMapTbl[]

Index	Bit Mask(Binary)
0	00000001
1	00000010
2	00000100
3	00001000
4	00010000
5	00100000
6	01000000
7	10000000

找出最高優先權的Task



```
y = OSUnMapTbl[pevent->OSEventGrp];  
x = OSUnMapTbl[pevent->OSEventTbl[y]];  
prio = (y << 3) + x;
```

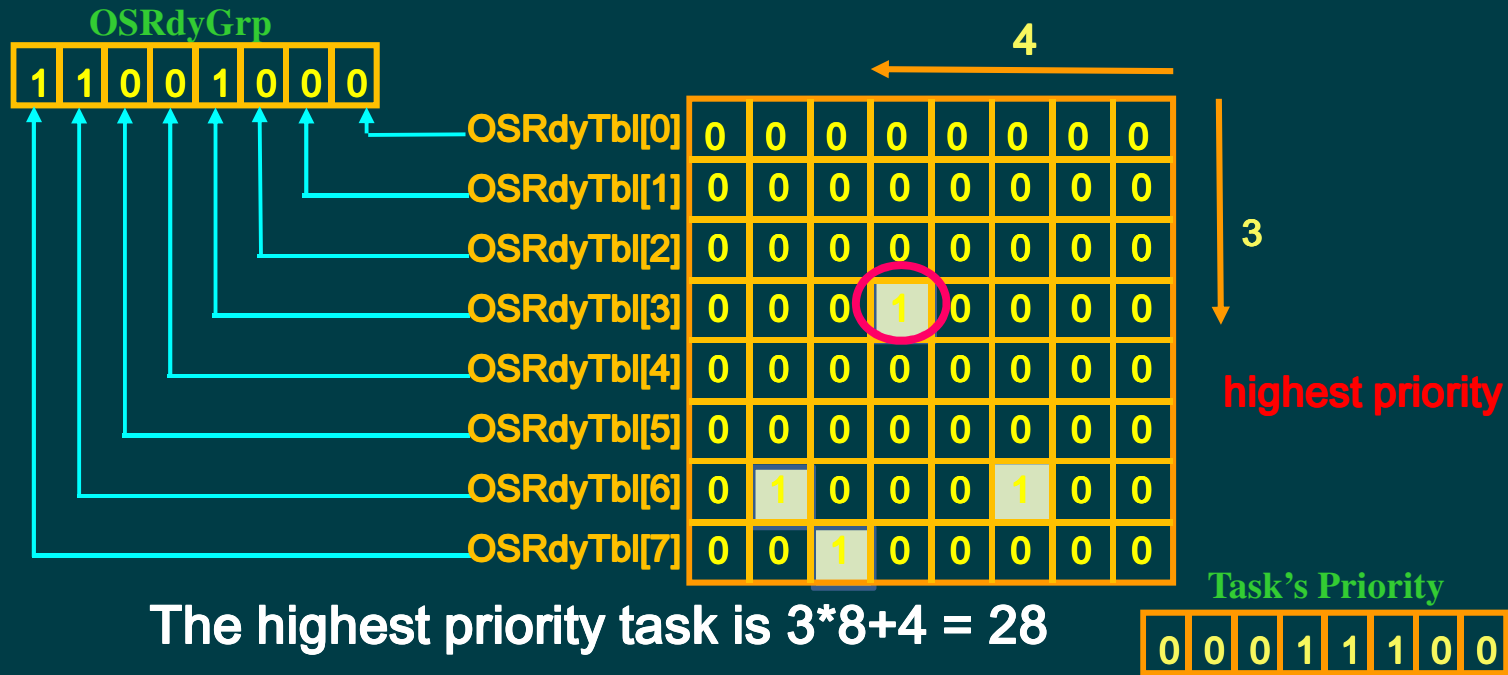
OSUnMapTbl[]

```
INT8U const OSUnMapTbl[] = {  
    0, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    6, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    7, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    6, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    5, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0,  
    4, 0, 1, 0, 2, 0, 1, 0, 3, 0, 1, 0, 2, 0, 1, 0};
```


Example: 由ECB wait list挑出最高優先權Task

OSEventGrp = 11001000 => OSUnMapTbl[OSEventGrp] = 3

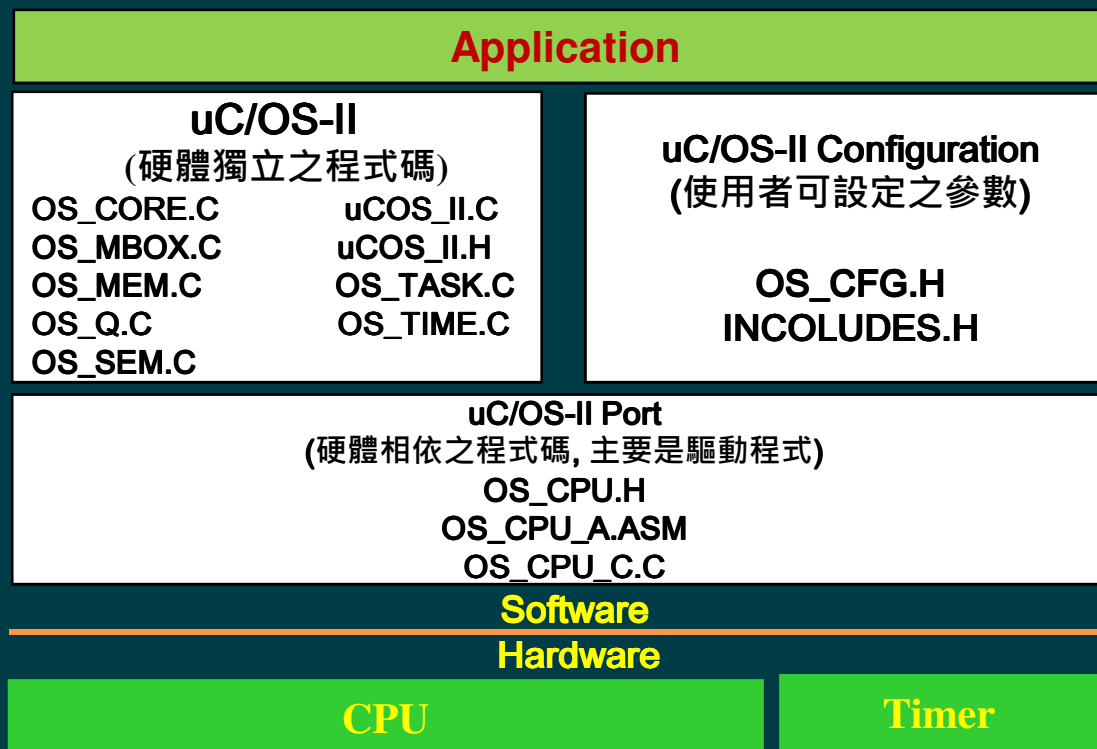
OSEventTbl[3] = 00010000 → OSUnMapTbl[OSEventTbl[3]] = 4



Part 3

uC/OS-II軟體架構

uC/OS-II軟/硬體架構



撰寫使用者程式

- 宣告task的stack
- 撰寫各別task之內容
- 撰寫main()程式碼(程式進入點)以對OS初始化與啟動
 - 對MCU, I/O周邊與OS做初始化
 - 啟動timer, 提供tick clock
 - 創建系統所需之semaphore, mail box, ...
 - 創建task
 - 啟動OS

uC/OS-II RTOS核心啟動

```
OS_STK TaskStartStk[TASK_STK_SIZE],
void main (void)
{
    OSInit(); //create Idle task and statistical task
    RandomSem = OSSemCreate(1);
    OSTaskCreate(TaskStart , //Task name
                (void *)0, //passing parameter
                (void *)&TaskStartStk[0], //Task stack
                0); //priority number
    OSStart(); //啟動作業系統
}
```

```
void TaskStart (void *p_arg)
{
    /* Initial System Tick */
    #if (OS_TASK_STAT_EN > 0)
        OSStatInit
    #endif
    /*Create your application tasks */
    for(;;) {
        /* Code for TaskStart() goes here! */
    }
}
```

撰寫Task內容

- 宣告Task所需之stack
 - 需為全域陣列變數，其變數類型為 OS_STK
- 所創建的task可為
 - 無窮迴圈的函數: 每一次迴圈必需呼叫下列函數以放棄CPU
 - OSSemPend(), OSMboxPend(), OSQPend(), OSTaskSuspend() (為事件觸發/同步Task)
 - OSTimeDly(), OSTimeDlyHMSM() (為周期性Task)
 - 只執行一次的Task: 函數的最後需執行OSTaskDel()，讓自己不再被scheduler排程

一般Task Structure

- 正常Task程式架構為下列兩種

```
void YourTask (void *pdata)
{
    for (;;) {
        /* USER CODE */
        Call one of uC/OS-II's services:
        OSMsgBoxPend(),
        OSQPend();
        OSSemPend();
        OSTaskDel(OS_PRIO_SELF);
        OSTaskSuspend(OS_PRIO_SELF);
        OSTimeDly();
        OSTimeDlyHMSM();
    }
} 等待事件Task
```

```
void YourTask (void *pdata)
{
    /* USER CODE */
    OSTaskDel(OS_PRIO_SELF);
} 只執行一次的Task
```

創建Task

- 可利用下列兩個函數來創建task:

1. OSTaskCreate()
2. OSTaskCreateExt()

- OSTaskCreate() 需要 4 參數:

1. 'tname' 為指標變數，指向要創建之task程式位址
2. 'pdata' 為指標變數，當task第一次執行時，指向提供給創建task之參數
3. 'ptos' 為指標變數，指向創建之stack啟始位址
3. 'prio' 所要創建task之優先權

```
INT8U OSTaskCreate(void (*task)(void *tname),  
void *pdata,  
OS_STK *ptos,  
INT8U prio);
```


Task創建範例

```
OS_STK Task1Stk[1024]; ←  
void main (void)  
{  
  INT8U err;  
  ...  
  OSInit(); /* Initialize µC/OS-II */  
  ...  
  OSTaskCreate(Task1,  
               (void *)0,  
               &Task1Stk[1023],  
               25);  
  ...  
  OSStart(); /* Start Multitasking */  
}
```

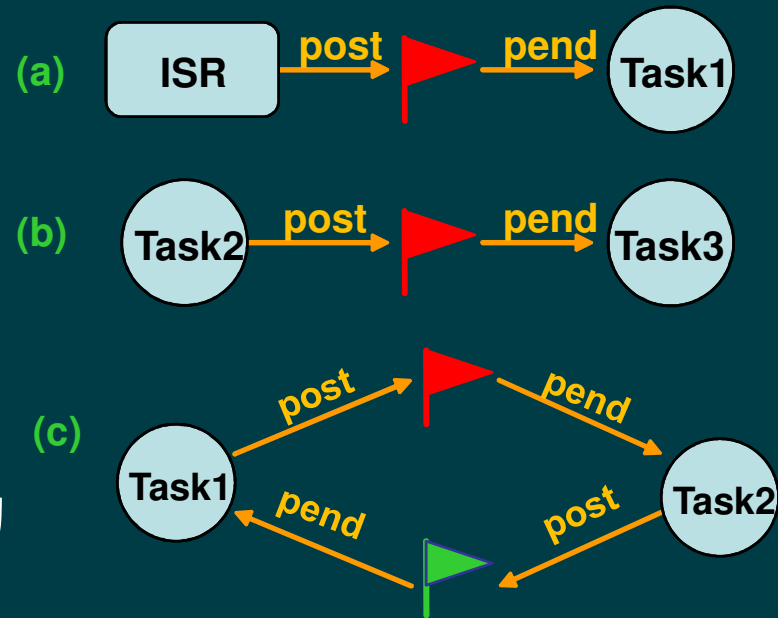
```
void Task1 (void *p_arg)  
{  
  (void)p_arg; /* Prevent compiler warning */  
  for (;;) {  
    ... /* Task code */  
    ...  
  }  
}
```

Part 4

Task同步-Semaphore

Task間同步-Semaphore

- 單一事件觸發的同步
 - ISR程式觸發Task
 - Task觸發Task
- 不可由ISR觸發ISR與Task觸發ISR
 - 會造成ISR程式執行時間不確定→影響系統即時性能力



uC/OS-II -Semaphore

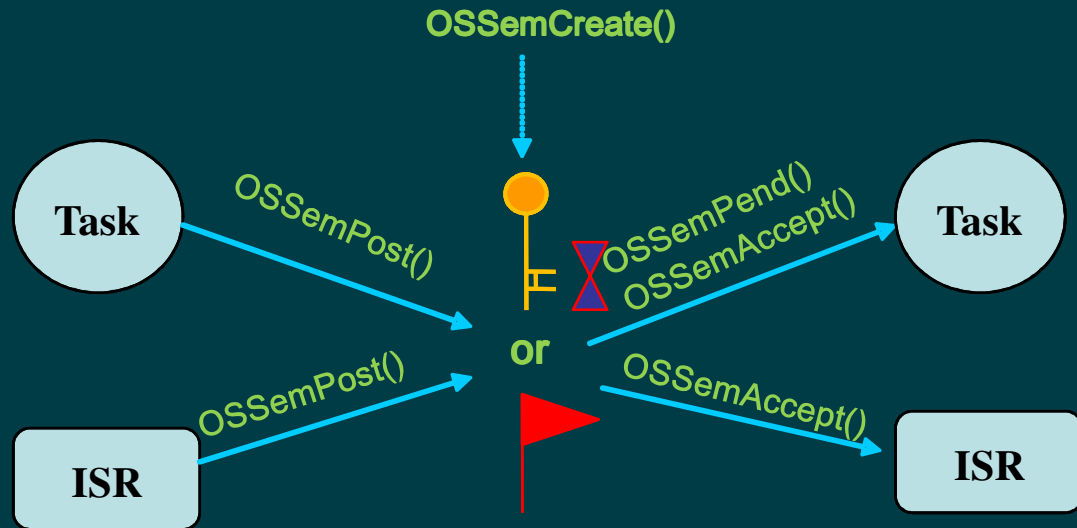
Semaphore相關之函數

➡ OSSemCreate()

OSSemPost()

OSSemPend()

OSSemAccept()



OSSemCreate() 函數範例

```
OS_EVENT *DispSem; ←  
  
void main (void)  
{  
  ...  
  OSInit(); /* Initialize µC/OS-II */  
  ...  
  DispSem = OSSemCreate(1); /* Create Display Semaphore */  
  ...  
  OSStart(); /* Start Multitasking */  
}
```

OSSemPost() 函數範例


```
OS_EVENT *DispSem; ←
void TaskX (void *p_arg)
{ INT8U err;
  (void)p_arg;
  for (;;) {
    ...
    err = OSEmPost(DispSem);
    switch (err) {
      case OS_ERR_NONE:
        /* Semaphore signaled */
        break;
      case OS_ERR_SEM_OVF:
        /* Semaphore has overflowed */
        break;
    }
  }
}
```

OSSemPend()函數範例

```
OS_EVENT *DispSem; ←
void DispTask (void *p_arg)
{
    INT8U err;
    (void)p_arg;
    for (;;) {
        ...
        OSSemPend(DispSem, 0, &err);
        ./* The only way this task continues is if _ */
        ./* _ the semaphore is signaled! */
    }
}
```

OSSemAccept() 函數範例

```
OS_EVENT *DispSem;
void DispTask (void *p_arg)
{
    INT8U err;
    (void)p_arg;
    for (;;) {
        ...
        value = OSSEMaccept(DispSem);
        if (value > 0) {
            ... /* Resource available, process */
        }
    }
}
```

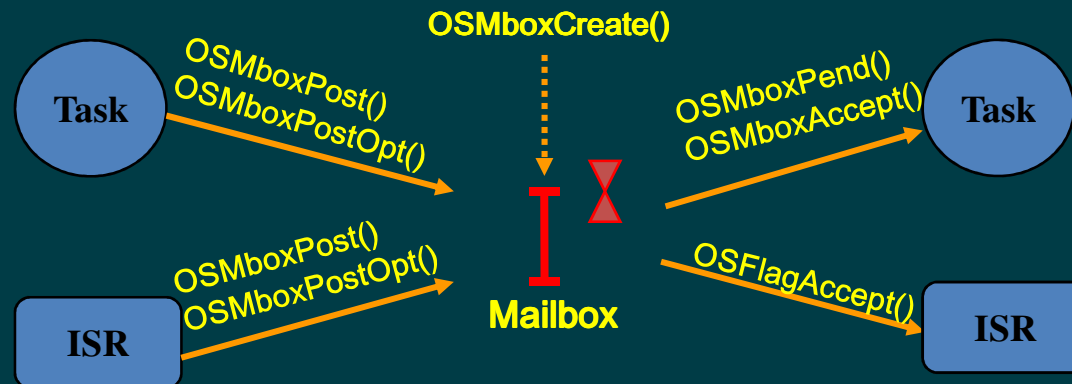


Part 5

uC/OS-II 訊息傳遞-Mailbox

Mailbox相關之函數

Mailbox相關之函數	OS_CFG.H需設定之參數
→ OSMboxCreate()	
OSMboxPost()	OS_Mbox_Post_EN = 1
OSMboxPostOpt()	OS_Mbox_Post_OPT_EN = 1
OSMboxPend()	
OSMboxAccept()	OS_Mbox_ACCEPT_EN = 1



OSMboxCreate() 函數範例

```
OS_EVENT *CommMbox; ←
void main (void)
{
    ...
    OSInit(); /* Initialize μC/OS-II */
    ...
    CommMbox = OSMboxCreate((void *)0); /* Create COMM mailbox */
    OSStart(); /* Start Multitasking */
}
```

OSMboxPost()函數範例

```
OS_EVENT *CommMbox; ←
INT8U CommRxBuf[100];


void CommTaskRx (void *p_arg)
{
    INT8U err;
    (void)p_arg;
    for (;;) {
        ...
        err = OSMboxPost(CommMbox, (void *)&CommRxBuf[0]);
        ...
    }
}
```

OSMboxPostOpt() 函數範例

```
OS_EVENT *CommMbox; ←
INT8U CommRxBuf[100];
void CommRxTask (void *p_arg)
{
    INT8U err;
    (void)p_arg;
    for (;;) {
        ...
        err = OSMboxPostOpt(CommMbox,
                           (void *)&CommRxBuf[0],
                           OS_POST_OPT_BROADCAST);
        ...
    }
}
```

OSMboxPend()函數範例

```
OS_EVENT *CommMbox;
void CommTask(void *p_arg)
{
    INT8U err;
    void *pmsg;
    (void)p_arg;
    for (;;) {
        ...
        pmsg = OSMboxPend(CommMbox, 10, &err);
        if (err == OS_ERR_NONE) {
            ...
            ./* Code for received message */
        } else {
            ...
            ./* Code for message not received within timeout */
        }
        ...
    }
}
```



OSMboxAccept() 函數範例

```
OS_EVENT *CommMbox;
void Task (void *p_arg)
{
    void *pmsg;
    (void)p_arg;
    for (;;) {
        pmsg = OSMboxAccept(CommMbox); /* Check mailbox for a message */
        if (pmsg != (void *)0) {
            ... /* Message received, process */
        } else {
            ... /* Message not received, do .. */
            ... /* .. something else */
        }
        ...
    }
}
```



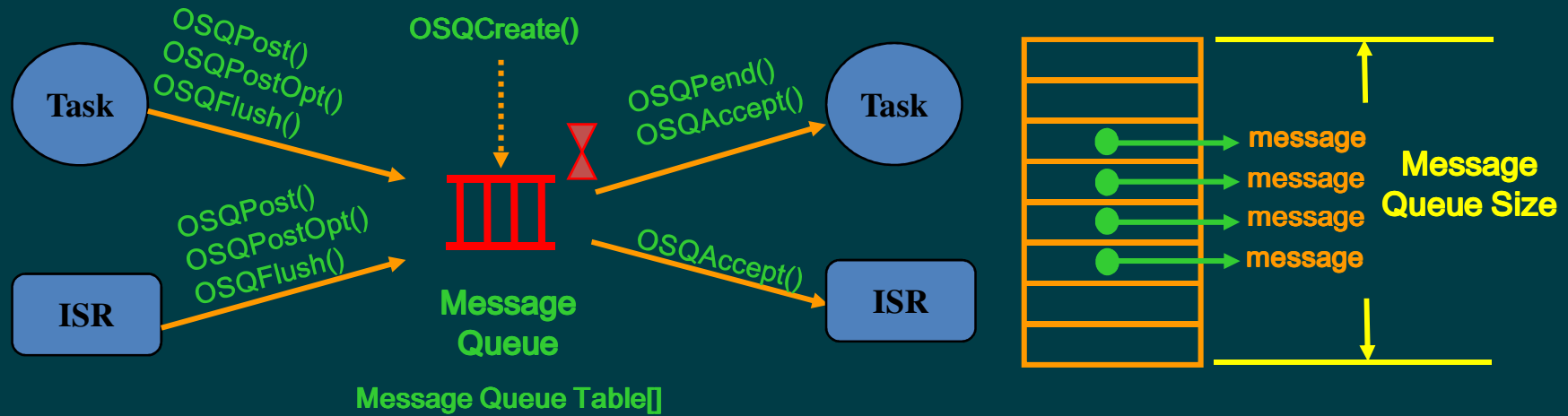
Part 6

Message Queue

Message Queue相關之函數

相關之函數	OS_CFG.H需設定之參數
→ OSQCreate()	
OSQPost()	OS_Q_Post_EN = 1
OSQPostOpt()	OS_Q_Post_OPT_EN = 1
OSQPend()	
OSQAccept()	OS_Q_ACCEPT_EN = 1
OSQFlush()	OS_Q_FLUSH_EN = 1

Message Queue



OSQCreate()函數範例

```
OS_EVENT *CommQ; ←
void *CommMsg[10];
void main (void)
{
    OSInit(); /* Initialize μC/OS-II */
    ...
    CommQ = OSQCreate(&CommMsg[0], 10); /* Create COMM Q */
    ...
    OSStart(); /* Start Multitasking */
}
```

Queue Head

Queue Length

OSQPost() 函數範例

```
OS_EVENT *CommQ;
INT8U CommRxBuf[100];
void CommTaskRx (void *p_arg)
{
    INT8U err;
    (void)p_arg;
    for (;;) {
        ...
        err = OSQPost(CommQ, (void *)&CommRxBuf[0]);
        switch (err) {
            case OS_ERR_NONE:
                /* Message was deposited into queue */
                break;
            case OS_ERR_Q_FULL:
                /* Queue is full */
                break;
            ...
        }
        ...
    }
}
```

←

Queue name

Message address

OSQPostOpt()函數範例

```
OS_EVENT *CommQ;  
INT8U CommRxBuf[100];
```



```
void CommRxTask (void *p_arg)  
{
```

```
    INT8U err;  
    (void)p_arg;  
    for (;;) {
```

```
        ...  
        err = OSQPostOpt(CommQ,
```

Queue name



Message address

```
        (void *)&CommRxBuf[0],  
        OS_POST_OPT_BROADCAST);
```



廣播此訊息給所有等待的Tasks

```
        ...  
    }
```

```
}
```

OSQPend()函數範例

```
OS_EVENT *CommQ;
void CommTask(void *p_arg)
{
    INT8U err;
    void *pmsg;
    for (;;) {
        ...
        pmsg = OSQPend(CommQ, 100, &err);
        if (err == OS_ERR_NONE) {
            ...
            /* Code for received message within 100 ticks*/
        } else {
            ...
            /* Code for message not received, must have timeout */
        }
        ...
    }
}
```

Queue name

最多等待100個ticks

OSQAccept() 函數範例

```
OS_EVENT *CommQ;
```

```
void Task (void *p_arg)  
{  
    void *pmsg;  
    (void)p_arg;  
    for (;;) {
```

```
        pmsg = OSQAccept(CommQ); /* Check message queue for a message */  
        if (pmsg != (void *)0) {  
            ... /* Message received, process */  
        } else {  
            ... /* Message not received, do .. */  
            ... /* .. something else */  
        }  
        ...  
    }  
}
```

Queue name



OSQFlush() 函數範例

```
OS_EVENT *CommQ;
```



```
void main (void)
```

```
{
```

```
    INT8U err;
```

```
    OSInit(); /* Initialize  $\mu$ C/OS-II */
```

```
    ...
```

```
    err = OSQFlush(CommQ);
```

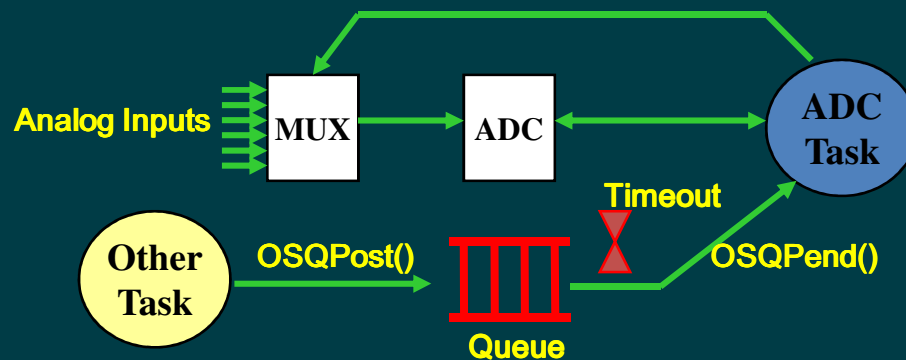
```
    ...
```

```
    OSStart(); /* Start Multitasking */
```

```
}
```


Message Queue於ADC環境之應用

- 利用Message Queue提供週期性的ADC轉換與即時轉換需求

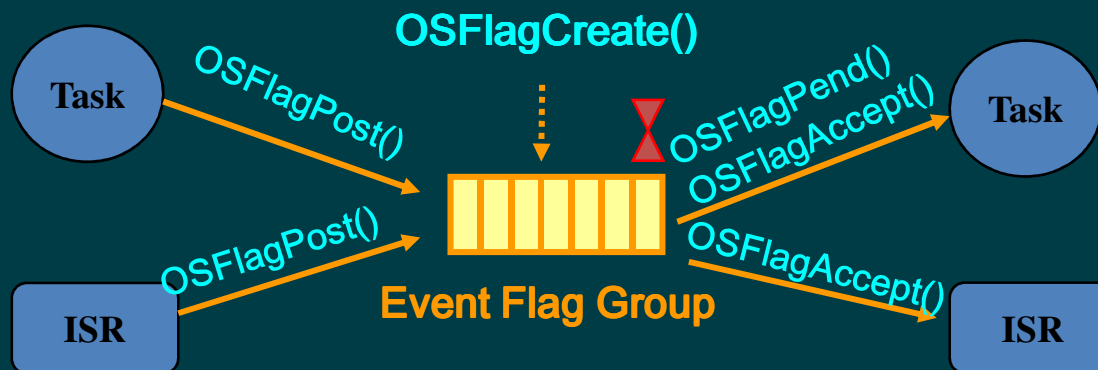


Part 7

uC/OS-II多事件觸發-Flags

Event Flags相關之函數

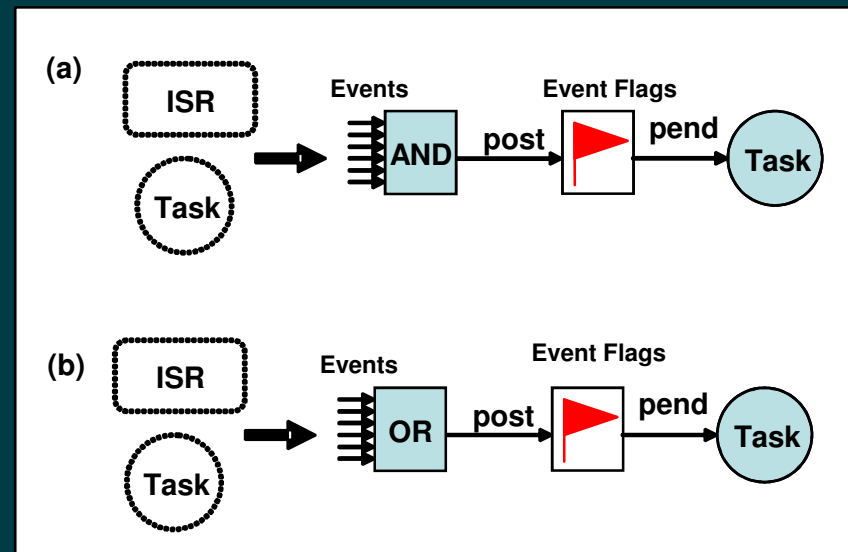
Event Flags相關之函數	OS_CFG.H需設定之參數
→ OSFlagCreate()	
OSFlagPost()	
OSFlagPend()	
OSFlagAccept()	OS_Flag_ACCEPT_EN = 1



Task間同步-Event Falgs

– 多事件觸發同步

- 多事件皆發生時觸發Task執行 (AND事件)
- 多事件中任一事件發生觸發Task執行(OR事件)



OSFlagCreate()

```
OS_FLAG_GRP *OSFlagCreate(OS_FLAGS flags,  
                           INT8U *perr);
```


- 參數:
 - **flags**: 為設定所產生的Event Flags初始值。
 - **perr**: 為一指標，指向錯誤碼變數，其可能之變數值為:
 - **OS_ERR_NONE**: 代表成功產生Event Flags。
- 回傳值:
 - OSFlagCreate()函數回傳值為一個指向Event Flags結構的指標，若已不允許產生新的Event Flags時，其回傳值為「NULL」

OSFlagCreate() 函數範例

```
OS_FLAG_GRP *EngineStatus; ←  
void main (void)  
{  
    INT8U err;  
    ...  
    OSInit(); /* Initialize μC/OS-II */  
    ...  
    /* Create a flag group containing the engine's status */  
    EngineStatus = OSFlagCreate(0x00, &err);  
    ...  
    OSStart(); /* Start Multitasking */  
}
```

OSFlagPost()

```
OS_FLAGS OSFlagPost(OS_FLAG_GRP *pgrp,  
                    OS_FLAGS flags,  
                    INT8U opt,  
                    INT8U *perr);
```




- 參數:
 - **pgrp**: 指向Event Flags結構的指標
 - **flags**: 描述要將Event Flags的那幾個位元設為1或0(由opt參數決定)，若opt參數為OS_FLAG_SET，則將flags參數所指定的位元設為1；若opt為OS_FLAG_CLR，則設為0。舉例來說，若想設定Event Flags的第0, 4, 5位元，則flag需設為0x31 (00110001)
 - **opt**: 描述 flags 所指定的位元要設為 1 (OS_FLAG_SET) or 0 (OS_FLAG_CLR).
 - **perr**: 指向執行結果之錯誤碼變數
- 回傳值: 回傳設定後的Event Flags值

OSFlagPost() 函數範例

```
#define ENGINE_OIL_PRES_OK 0x01 ←
#define ENGINE_OIL_TEMP_OK 0x02
#define ENGINE_START 0x04
OS_FLAG_GRP *EngineStatusFlags;
void TaskX (void *p_arg)
{
    INT8U err;
    (void)p_arg;
    for (;;) {
        ...
        err = OSFlagPost(EngineStatusFlags,
                        ENGINE_START,
                        OS_FLAG_SET,
                        &err);
        ...
    }
}
```


OSFlagPend()

```
OS_FLAGS OSFlagPend(OS_FLAG_GRP *pgrp,  
                    OS_FLAGS flags,  
                    INT8U wait_type,  
                    INT32U timeout,  
                    INT8U *perr);
```



- **pgrp**: 為指向某Event Flags結構的指標。
- **flags**: 設定所要等待Event Flags的那些事件位元。
- **timeout**: 設定要等待Event Flags事件之時間，若等待時間Timeout會返回，但會回傳一個錯誤碼，當此值設為「0」時，表無條件等待直至事件位元組發生。
- **perr**: 表執行此函數是否正確，若其值為**OS_ERR_NONE**，表正確執行此函數
- **回傳值**: 為0代表有問題，若非0值，其值代表所要等待Event Flags之值

OSFlagPend()


```
OS_FLAGS OSFlagPend(OS_FLAG_GRP *pgrp,  
                    OS_FLAGS flags,  
                    INT8U wait_type, ←  
                    INT32U timeout,  
                    INT8U *perr);
```

- **wait_type**: 決定是要等待Event Flags之位元事件為1或為0，其設定參數為：
 - **OS_FLAG_WAIT_CLR_ALL**: Event Flags的所有事件位元需皆為0
 - **OS_FLAG_WAIT_CLR_ANY**: 依**flags**參數所設定的事件位元皆為0
 - **OS_FLAG_WAIT_SET_ALL**: Event Flags的所有事件位元需皆為1
 - **OS_FLAG_WAIT_SET_ANY**: 依**flags**參數所設定的事件位元皆為1
- 若要將所發生的事件位元清為0，則此**wait_type**參數可設為
 - **OS_FLAG_WAIT_SET_ANY + OS_FLAG_CONSUME**

OSFlagPend() 函數範例

```
#define ENGINE_OIL_PRES_OK 0x01
#define ENGINE_OIL_TEMP_OK 0x02
#define ENGINE_START 0x04
OS_FLAG_GRP *EngineStatus;
void Task (void *p_arg)
{
    INT8U err;
    OS_FLAGS value;
    (void)p_arg;
    for (;;) { value = OSFlagPend(EngineStatus,
        ENGINE_OIL_PRES_OK + ENGINE_OIL_TEMP_OK,
        OS_FLAG_WAIT_SET_ALL + OS_FLAG_CONSUME,
        10,
        &err);
```

OSFlagPend() 函數範例

```
switch (err) {  
    case OS_ERR_NONE:  /* Desired flags are available */  
        break;  
    case OS_ERR_TIMEOUT:  
        /* The desired flags were NOT available before .. */  
        /* .. 10 ticks occurred */  
        break;  
}  
...  
}
```

OSFlagAccept()

```
OS_FLAGS OSFlagAccept(OS_FLAG_GRP *pgrp,  
                      OS_FLAGS flags,  
                      INT8U wait_type,  
                      INT8U *perr);
```


- 參數：其參數值與OSFlagPend()函數除了**timeout**參數外，其餘皆相同。

OSFlagAccept() 函數範例

```
#define ENGINE_OIL_PRES_OK 0x01 ←
#define ENGINE_OIL_TEMP_OK 0x02
#define ENGINE_START 0x04
OS_FLAG_GRP *EngineStatus;

void Task (void *p_arg)
{
    INT8U err;
    OS_FLAGS value;
    (void)p_arg;
    for (;;) { value = OSFlagAccept(EngineStatus,
                                   ENGINE_OIL_PRES_OK + ENGINE_OIL_TEMP_OK,
                                   OS_FLAG_WAIT_SET_ALL,
                                   &err);
```

OSFlagAccept() 函數範例

```
switch (err) {  
    case OS_ERR_NONE:  /* Desired flags are available */  
        break;  
    case OS_ERR_FLAG_NOT_RDY:  
        /* The desired flags are NOT available */  
        break;  
    ...  
}
```

Part 8

uC/OS-II記憶體管理

記憶體管理相關之函數


相關之函數	OS_CFG.H需設定之參數
→ OSMemCreate()	
OSMemGet()	
OSMemPut()	

- ◆ 是採用**Partition**方式的記憶體管理
 - 每個Partition內為**多個相同大小**的記憶體區塊
 - 利用**OSMemCreate()**函數設定對**所宣告的記憶體Partition**管理

OSMemCreate () 函數範例

```
OS_MEM *CommMem;
INT32U CommBuf[16][32];

void main (void)
{
    INT8U err;
    OSInit(); /* Initialize μC/OS-II */
    ...
    CommMem = OSMemCreate(&CommBuf[0][0], 16, 32 * sizeof(INT32U), &err);
    ...
    OSStart(); /* Start Multitasking */
}
```

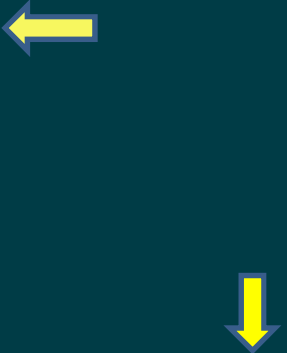


OSMemGet ()函數範例

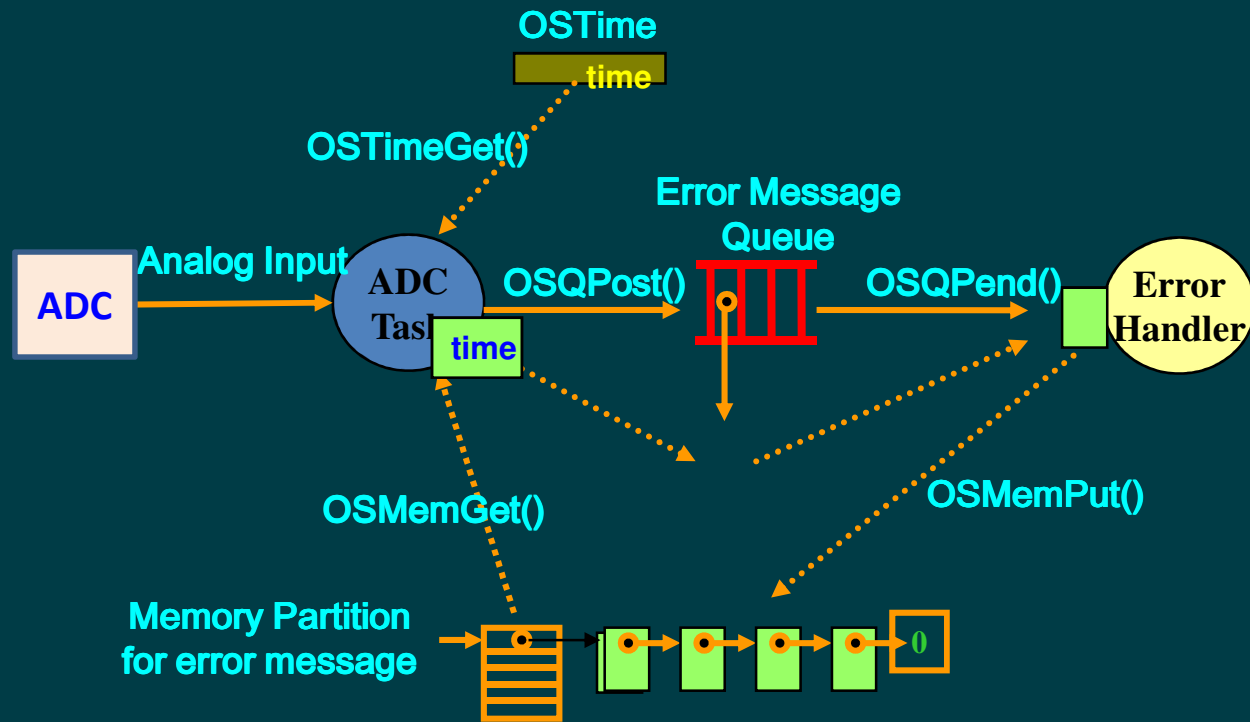
```
OS_MEM *CommMem; ←  
void Task (void *p_arg)  
{  
    INT8U *pmsg;  
    (void)p_arg;  
    for (;;) {  
        pmsg = OSMemGet(CommMem, &err);  
        if (pmsg != (INT8U *)0) {  
            ... /* Memory block allocated, use it. */  
        }  
        ...  
    }  
}
```

OSMemPut () 函數範例

```
OS_MEM *CommMem; ←
INT8U *CommMsg;
void Task (void *p_arg)
{
    INT8U err;
    (void)p_arg;
    for (;;) {
        err = OSMemPut(CommMem, (void *)CommMsg);
        if (err == OS_ERR_NONE) {
            .... /* Memory block released */
        }
        ...
    }
}
```



Memory Partition應用



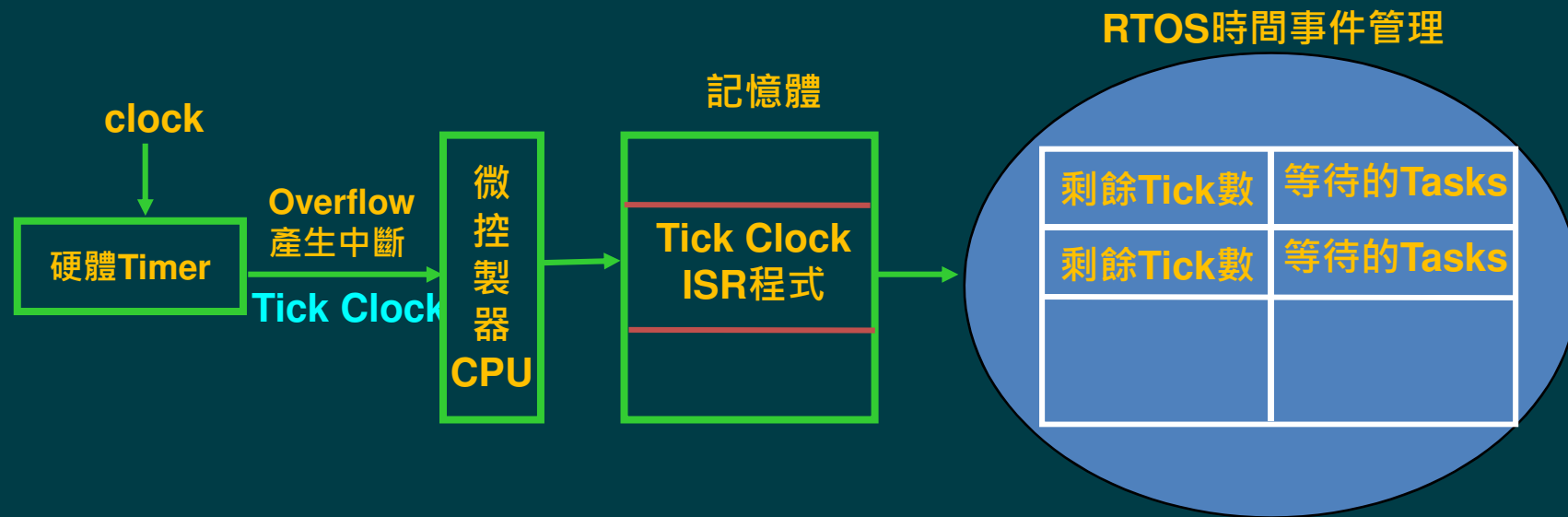
Part 9

uC/OS-II時間管理

RTOS為何需要Timer?

- Round-Robin排程器，每個Task執行一固定slice time即context switch，換下一個Task執行
- Task要周期性的polling周邊裝置之狀況
- Task在設定時間內，等待某事件發生，若時間已到，事件尚未發生，則Task將不繼續等待

RTOS時間事件管理



時間管理

- **Tick Clock**為RTOS一個重要的時間參考依據
 - 為一週期性的中斷(Tick)

時間管理相關之函數

相關之函數	OS_CFG.H需設定之參數
OSTimeDly()	
→ OSTimeDlyHMSM()	OS_Q_Post_EN = 1
OSTimeGet()	OS_Q_Post_OPT_EN = 1
OSTimeSet()	

OSTimeDly ()函數範例

```
void TaskX (void *p_arg)
{
for (;;) {
    ...
    OSTimeDly(10); /* Delay task for 10 clock ticks */
    ...
}
}
```

OSTimeDlyHMSM ()函數範例

```
void TaskX (void *p_arg)
{
for (;;) {
    ...
    OSTimeDlyHMSM(0, 0, 1, 0); /* Delay task for 1 second */
    ...
}
}
```

OSTimeGet () 函數範例

```
void TaskX (void *p_arg)
{
  INT32U clk;
  for (;;) {
    ...
    clk = OSTimeGet();
    ...           /* Get current value of system clock */
  }
}
```

OSTimeSet () 函數範例

```
void TaskX (void *p_arg)
{
for (;;) {
    ...
    OSTimeSet(0L); /* Reset the system clock */
    ...
}
}
```